



# A semantics for React

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# ReactFX: reactive programming with effect handlers

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# ReactFX: reactive programming with effects and handlers

## Research objectives

- **Foundations:** unify synchronous effects from the programming language with asynchronous events from the environment
- **Implementations:** experiment with research languages, e.g., Links, Koka, OCaml
- **Case studies:** React-style web applications, spreadsheets, etc.
- **Effect typing:** exploit for optimisation and modularity
- **Incremental updates:** use effect handlers to abstract over incremental updates to virtual DOM
- **Pre-emptive concurrency:** synergy with Cangjie implementation of effect handlers

# Understanding React

- What are components?
  - They look like stateful functions — *in what sense are they functions?*
- Goal: **translate** React into something more well-understood
  - **Expressivity** result [Felleisen, 1991]
  - What programming language features are required?
- We define the following:
  - $\lambda_{\text{react}}$  : a simplified React-like source language (inspired by React-tRace [Lee, Ahn, Yee, OOPSLA 2025])
  - $\lambda_{\text{ref}}$  : a canonical target language with ML-style references
  - a **translation**  $\llbracket - \rrbracket$  from  $\lambda_{\text{react}}$  into  $\lambda_{\text{ref}}$

# Core language

Values  $v, w ::= x \mid () \mid s \mid \lambda x.e \mid [] \mid v :: w$

$\mid \text{text } v \mid \text{tag}(v, w) \mid \text{attr}(v, w) \mid \text{onClick } v$

Computations  $e ::= \text{return } v \mid \text{let } x = e \text{ in } e' \mid v w$

$\mid \text{case } v \{ [] \Rightarrow e \mid x :: xs \Rightarrow e' \}$

Types  $A, B ::= \mathbf{1} \mid \text{String} \mid A \rightarrow B \mid \text{List } A \mid \text{View} \mid \text{Attr}$

- Fine-grain call-by-value
- Types & constructors for HTML

# Source language: $\lambda_{\text{react}}$

Values  $v, w ::= \dots \mid \text{comp}\{b\} \mid \text{set}@_p^\ell$

Bodies  $b ::= e \mid \text{let } x = e \text{ in } b \mid \text{let } x, x_{\text{set}} = \text{useState}^\ell v \text{ in } b$

Types  $A, B ::= \dots \mid A \xrightarrow{E} B$

Effect  $E ::= \text{Pure} \mid \text{Mut}$

Record  $\mathcal{R} ::= \emptyset \mid \mathcal{R}, \ell : A$

Syntactically force useState to occur only at component top-level!

# React versus $\lambda_{\text{react}}$

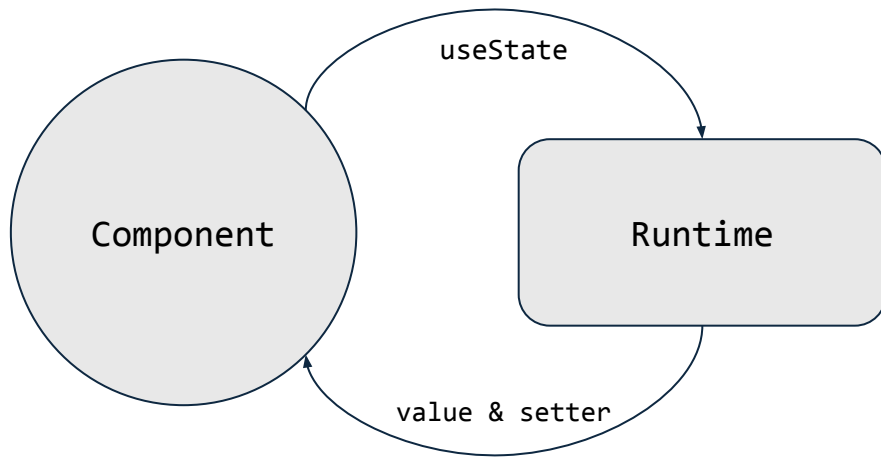
```
const App = z => {  
  const [x, x_set] = useState(z);  
  const y = x * x;  
  return (  
    <p onClick={() => {x_set (x + 1)}}>  
      {y}  
    </p>  
  )  
}
```

```
let App = λz.comp {  
  let x, x_set = useState z in  
  let y = x * x in  
  return (  
    <p onClick={λ().{x_set (x + 1)}}>  
      {y}  
    </p>  
  )  
}
```

# $\lambda_{\text{react}}$ design

- Looks & acts like React
- **Abstract away** from details
- Focus on “essence” of React
  - Components
  - State
  - Runtime

```
> ('b' + 'a' + + 'a' + 'a').toLowerCase()  
< 'banana'
```



# $\lambda_{\text{react}}$ is type-safe!

```
const App = z => {  
  const [x, x_set] = useState(z);  
  x_set (x + 1);  
  const y = x * x;  
  return (  
    <p onClick={() => {x_set (x + 1)}}>  
      {y}  
    </p>  
  )  
}
```

—————> Infinite re-render!

# $\lambda_{\text{react}}$ is type-safe!

```
const App = z => {
  const [x, x_set] = useState(z);
  x_set (x + 1);
  const y = x * x;
  return (
    <p onClick={() => {x_set (x + 1)}}>
      {y}
    </p>
  )
}
```

```
let App = λz.comp {
  let x, x_set = useState z in
  let () = x_set (x + 1) in
  let y = x * x in
  return (
    <p onClick={λ().{x_set (x + 1)}}>
      {y}
    </p>
  )
}
```

# $\lambda_{\text{react}}$ is type-safe!

Error: Mutation effect  
introduced in component body!

$$\frac{\Gamma \vdash v_1 : A \xrightarrow{\text{Mut}} \mathbf{1} \quad \Gamma \vdash v_2 : A}{\Gamma \vdash v_1 v_2 : \mathbf{1} ! \text{Mut}} \text{TE-APP}$$

$$\frac{\Gamma \vdash e : A ! \text{Pure} \quad \Gamma, x : A \vdash \hat{e} : \mathcal{R}}{\Gamma \vdash \text{let } x = e \text{ in } \hat{e} : \mathcal{R}} \text{TB-LET}$$

```
let App =  $\lambda z$ .comp {  
  let x, xset = useState z in  
  let () = xset (x + 1) in  
  let y = x * x in  
  return (  
    <p onClick={ $\lambda()$ .{xset (x + 1)}}>  
      {y}  
    </p>  
  )  
}
```

Target language:  $\lambda_{\text{ref}}$

Values  $v, w ::= \dots \mid r$

Computations  $e ::= \dots \mid \text{ref } v \mid !v \mid v := w$

Types  $A, B ::= \dots \mid \text{Ref } A$

# $\lambda_{\text{react}}$ to $\lambda_{\text{ref}}$ by example

```
 $\lambda z.$ comp {  
  let x, xset = useState z in  
  let y = x * x in  
  return (  
    <p onClick={ $\lambda()$ .{xset (x + 1)}}>  
      {y}  
    </p>  
  )  
}
```

$\lambda z.$  $\lambda()$ .

```
let x = return z in  
let x1 = ref(x) in  
let xset =  $\lambda w.$ x1 := w in  
return  $\lambda()$ .
```

Initial

```
let x = !x1 in  
let xset =  $\lambda w.$ x1 := w in
```

Successive

```
let y = x * x in  
return (  
  <p onClick={ $\lambda()$ .{xset (x + 1)}}>  
    {y}  
  </p>  
)
```

$\lambda_{\text{react}}$  to  $\lambda_{\text{ref}}$

- Key insight: components have *two* different interpretations!
  - When *first* run, useState **initializes** a reference
  - On *successive* runs, useState **retrieves** that reference

$$\llbracket \text{let } x, x_{\text{set}} = \text{useState}^{\ell} v \text{ in } b \rrbracket^{\text{Init}} \approx \text{let } x = \text{return } \llbracket v \rrbracket \text{ in}$$
$$\quad \text{let } x_{\ell} = \text{ref}(x) \text{ in}$$
$$\quad \text{let } x_{\text{set}} = \lambda y. x_{\ell} := y \text{ in}$$
$$\quad \llbracket b \rrbracket^{\text{Init}}$$

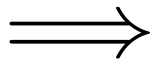
$\lambda_{\text{react}}$  to  $\lambda_{\text{ref}}$

- Key insight: components have *two* different interpretations!
  - When *first* run, useState **initializes** a reference
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$$\llbracket \text{let } x, x_{\text{set}} = \text{useState}^{\ell} v \text{ in } b \rrbracket^{\text{Succ}} \approx \text{let } x = !x_{\ell} \text{ in}$$
$$\text{let } x_{\text{set}} = \lambda y. x_{\ell} := y \text{ in}$$
$$\llbracket b \rrbracket^{\text{Succ}}$$

# What next for $\lambda_{\text{react}}$ ?

$$m, e \mapsto e', m'$$



$$\llbracket m \rrbracket, \llbracket e \rrbracket \mapsto^* \llbracket e' \rrbracket, \llbracket m' \rrbracket$$

- Proof of operational correspondence
  - Show our translation is “good” (almost done)

- Possible extensions: capture more of React
  - useEffect?
  - Reconciliation?
  - Custom hooks

# Future directions: beyond $\lambda_{\text{react}}$

- How do we **react** with the outside world?
  - $\lambda_{\text{react}}$  assumes discrete, synchronous events
  - What about continuous signals?
    - Draw on **functional reactive programming**
- How do effect handlers fit into this picture?
  - Connecting effect handlers and React
  - Asynchronous effects [[Ahman and Pretnar, 2024](#)]
  - ...